

- Cost enum works if $|G:H|$ finite,
but this is unsolvable (so procedures $\rightarrow$ not alg).
- Can do: compute all subgrps of index $\leq n$. (fixed)
 finding pres if order & mult table known.

Presentations of subgroups: theory

Recall: $G = \langle X \mid R \rangle$, and $G \cong F/N$ where $F = F(X)$ and $N = \langle\langle R \rangle\rangle$.
 Let $H = E/N < G$ and let $\tilde{T} \subseteq F$ be a right transversal of E in F that maps onto the right transversal T of H in G :

$$F = \bigcup_{\tilde{t} \in \tilde{T}} E\tilde{t} \quad \text{and thus} \quad G = \bigcup_{t \in T} Ht$$

We assume $1_F \in \tilde{T}$. For $w \in F$, define $\bar{w} := t \in T$ with $w \in Et$.

Thm (Nielsen - Schreier)

$E = \langle Y \rangle$ and hence $H = \langle yN : y \in Y \rangle$, where
 $Y := \{ t_x (\overline{t_x})^{-1} \mid t \in T, x \in X, t_x \neq \bar{t}_x \}$
 $\subseteq F \setminus \{1_F\}$.

If T is prefix-closed, then E is freely
 gen by Y .

Thm (Reide - Schreier)

Let G, F, H, E, T above & prefix closed.

Then

$$H \cong \langle Y \mid p(twt^{-1}) \forall t \in T, w \in R \rangle$$

$P: E \rightarrow F(Y)$ is isom mentioned above.

Proof (ex)

Example (The symmetric group S_4)

Let $G := \langle x, y \mid x^3, y^4, (xy)^2 \rangle$, and $H := \langle x, yx^{-1}y^{-2} \rangle$. First, do coset enumeration. Now the coset table is

	x	y	x^{-1}	y^{-1}
1	1	2	1	3
2	3	4	4	1
3	4	1	2	4
4	2	3	3	2

underlined
= def of coset

$$1 = 1$$

$$2 = y$$

$$3 = yx$$

$$4 = y^2$$

Example (The symmetric group S_4 (ctd))

Now, we can define the elements of the free generating set Y of E by inserting new symbols into the remaining entries on the first two columns of the coset table.

	x	y	x^{-1}	y^{-1}
1	a1	2	$a^{-1}1$	$c^{-1}3$
2	3	4	$d^{-1}4$	1
3	b4	c1	2	$e^{-1}4$
4	d2	e3	$b^{-1}3$	2

$$ex: b = (yx)x(y^{-2})$$

a, b, c, d = Schreier generators,

$$Y = \{a, b, c, d, e\} \text{ w/ } a = x, b = yx^2y^{-2}, c = yxy, \\ d = y^2xy^{-1}, e = y^3x^{-1}y^{-1}$$

Now compute $\rho(twt^{-1})$'s.

Example (The symmetric group S_4 (ctd))

We calculate the relators $\rho(twt^{-1})$ of H by scanning the cosets of H in G under the relators of G using the modified coset table.

	x		x		x	
1		a1		a ² 1		a ³ 1
2		3		b4		bd2

$$yx^3 = bd y$$

$$yx^3 y^{-1} = bd.$$

	y		y		y		y	
1		2		4		e3		ec1

	x		y		x		y	
1		a1		a2		a3		ac1
3		b4		be5		beb5		(be) ² 2
4		d2		d4		d ² 2		d ² 4

Now get presentation...

Example (The symmetric group S_4 (ctd))

So

$$H \cong \langle a, b, c, d, e \mid a^3, bd, ec, ac, (be)^2, d^2 \rangle$$

We can now perform Tietze transformations on this presentation and use the relators bd , ec , ac to eliminate $d = b^{-1}$, $e = c^{-1}$, $c = a^{-1}$, to give the presentation

$$H \cong \langle a, b \mid a^3, (ba)^2, b^2, \rangle$$

which is a standard presentation of the dihedral group of order 6 (and also of the symmetric group S_3).

Q : Can you do Tietze while running alg?

A : Yes. Ppl have tried.

- Can also find pres w/ user-supplied pres, but have to calc subgrp pres, while doing coset enum. If index large, can result in very long reductors.
- From D. Holt's exp: user-supplied $\sim$ same as CPU gen.
 - order of operations make big difference.
- Have to deal w/ coincidences (of course).

PROBLEM

Given $G = \langle x | R \rangle$.

How to prove $|G| = \infty$?

(for finite, can do coset enum).

SOL: 1) Find abelian invariants $\rightarrow$ that's easy
↳ maybe one of them is ∞ .

2) • Compute low-index subgrps

- Use Reid-Schreier to find pres.

- Compute ab. inv for those.

- If find infinite factor, Then $G \infty$ also.

Example

$$G := \langle x, y \mid x^3, y^3, (xy)^3 \rangle; \quad H := \langle x^{-1}y, yx^{-1} \rangle.$$

Show that $|G : H| = 3$, calculate a presentation of H and show that H has abelian invariants $[0, 0]$. (In fact $H \cong \mathbb{Z}^2$.)

EVERYTHING SO FAR

Computing in various quotients of fp
grps $G = \langle X \mid R \rangle$.

Next

Solving Dehn in specific grp classes.

I. Small Cancellation Th.

Small Canc. $\sim$ two rels don't have overlapping rels.

Dehn: on surface grp w/ small.canc.
prop.

Small cancellation theory

Let $G := \langle X \mid R \rangle$, and assume that all $r \in R$ are cyclically reduced. Let $\hat{R}$ be the set of all rotations and inversions of the elements of R .

For example, if $R = \{a^3, b^3, aba^{-1}b^{-1}\}$, then (with $\bar{a} := a^{-1}$, $\bar{b} := b^{-1}$):

$$\hat{R} = \{a^3, \bar{a}^3, b^3, \bar{b}^3, ab\bar{a}\bar{b}, b\bar{a}\bar{b}a, \bar{a}\bar{b}ab, \bar{b}ab\bar{a}, ba\bar{b}\bar{a}, \bar{a}\bar{b}a\bar{b}, \bar{b}\bar{a}ba, \bar{a}b\bar{a}\bar{b}\}.$$

clear
R

Def : (piece)

non-empty word that is prefix
of two dif elements of $\hat{R}$.

ie $pu, pv \in \hat{R}$ $u \neq v$.

Def ($C'(\lambda)$)

Small cancellation theory (ctd)

Example

$$G = \langle a, b, c, d \mid a^7, b^7, c^7, d^7, a^{-1}b^{-1}abc^{-1}d^{-1}cd \rangle.$$

The pieces are $a^{\pm 1}, b^{\pm 1}, c^{\pm 1}, d^{\pm 1}$, and G satisfies $C'(1/6)$.

pieces = gens & inv.

every fraction of word piece occupies
 $< 1/7$.

so $\lambda > 1/6$.

Definition (Condition $T(q)$)

We say that $\langle X \mid R \rangle$ is $T(q)$, if the following holds:

- Let $3 \leq h < q$ and $(r_1, r_2, \dots, r_h) \in \hat{R}^h$ with no successive elements r_i, r_{i+1} or r_h, r_1 an inverse pair. Then at least one of the products $r_1 r_2, r_2 r_3, \dots, r_h r_1$ is reduced without cancellation.

$T(q)$ says that all vertices in any reduced van Kampen diagram for the presentation have valency at least q .

Thm (Lyndon - Schupp)

(use curvature of VK-diag to prove results)

$G = \langle x \mid R \rangle$, $x \in R$ cyclically red,

if : $C'(C'10)$ or

$C'(14)$ and $T(4)$ or

$C'(13)$ and $T(3)$

then : Dehn's alg solves WP
for $G = \langle x \mid R \rangle$ in linear time.

Keep ex :

$$G = \langle a, b, c, d \mid a^7, b^7, c^7, d^7, a^{-1}b^{-1}abc^{-1}d^{-1}cd \rangle$$

Dehn alg : use Rewrite sys

Rewrite sys (RW)

Sub one str for another.

Rewrite systems

Rewrite systems arise in a wide variety of situations, such as text editing and formal language theory as well as group theory!

Definition (Rewrite system)

Let A be a finite set, known as the **alphabet**. A **rewrite system** (RWS) $\mathcal{R}$ over A is a set of **rules** $u \rightarrow v$ with $u, v \in A^*$.

We apply a rule $u \rightarrow v$ in $\mathcal{R}$ to a word $w \in A^*$ by replacing an occurrence of u as a subword of w by v . We then write $w \rightarrow_{\mathcal{R}} w'$ (or usually just $w \rightarrow w'$) where w' is the result of the replacement.

In other words, we write $svt \rightarrow swt$ for all $s, t \in A^*$.

We write $w \Rightarrow_{\mathcal{R}} w'$ if either $w = w'$, or w' results from w by a sequence of applications of rewrite rules from $\mathcal{R}$. That is: $w = w'$ or there is a finite tuple $(w_1, w_2, \dots, w_k)$ of words with

$$w \rightarrow w_1 \rightarrow w_2 \rightarrow \dots \rightarrow w_k \rightarrow w'$$

Note: $w \Rightarrow_R w'$ need NOT be from finite steps of rewrites.

Def (Dehn RW sys)

Given fp, $G = \langle X | R \rangle$

write all $r \in \hat{R}$ as $r = abu$ w/

$|a| > |b| > |a| - 2$ & define

RW $a \rightarrow b^{-1}$ over $A = X \cup X^{-1}$.

Dehn RWS = set of rules w/ $aa^{-1} \rightarrow 1$
 $\forall a \in A$.

Ex: $abcde \in R$

$$\downarrow$$
$$abc = e^{-1}d^{-1}$$

$$\Rightarrow abc \rightarrow e^{-1}d^{-1}$$

Dehn's alg solves WP

if $\forall w: w =_G 1 \iff w \Rightarrow_R 1$.

↑
irrespective
of the order
of rules.

"confluent on id"

- RW rules length red $\rightarrow$ solve WP in linear time.
- Works for when Grp satisfies some small cancellation property.

Dehn's algorithm: example

! (Small error somewhere)

Consider the word $w = a^3 b^6 a b c d^{-1} c d^{-6} a^4$ in the example above. Then

$$a^{-1} b^{-1} a b c^{-1} d^{-1} c d \in R \Rightarrow a b c^{-1} d^{-1} c d a^{-1} b^{-1} \in \hat{R}$$

so there is a rewrite rule $a b c^{-1} d^{-1} c \rightarrow b a^{-1} d^{-1}$.

Applying this to w gives

$$w = a^3 b^6 \underline{a b c d^{-1} c} d^{-6} a^4 \rightarrow a^3 b^6 \underline{b a^{-1} d^{-1}} d^{-6} a^4 = a^3 b^7 d^{-7} a^4.$$

Now $b^7, d^{-7}, a^7 \in \hat{R}$ yield rewrite rules $b^4 \rightarrow b^{-3}$, $d^{-4} \rightarrow d^3$, $a^4 \rightarrow a^{-3}$, and applying these to w , followed by applications of $b^{-1} b \rightarrow 1$, etc, gives:

$$a^3 b^7 d^{-7} a^4 \rightarrow a^3 b^{-3} b^3 d^3 d^{-3} a^4 \rightarrow a^7 \rightarrow a^{-3} a^3 \rightarrow 1$$

so $w =_G 1$.

Dehn WP solv lin $\leftrightarrow$ hyp. grp.

Q: What are the small cancellation properties of hyperbolic grps?

Theorem

The following conditions are equivalent for groups $G = \langle X \rangle$ with X finite.

- There is a set of relators $R \subset A^*$ such that $G = \langle X \mid R \rangle$ and the Dehn algorithm of this presentation solves the word problem for G ;
- The **Dehn function** of G (w.r.t any set of defining relators) is linear;
- Geodesic triangles in the Cayley graph $\Gamma(G, A)$ of G are uniformly slim (or uniformly thin);
- various other conditions.

Idea?



does save some time?

GAP Packages!

- Given fp hyp grp, calculate set of RW rules for Dehn alg

KEY WORDS: R. Parker, C. Roney-Davgal

- KBMA G $\rightarrow$ can compute automatic grps & check if hyp (slower)

Q $O(WP(G)) = n^2$ which G 's?

A:
• Automatic grps & others.
• Cox $\sim n^{1+\epsilon}$.

NOTE: • October $\rightarrow$ lectures ext next talk

• Nov $\rightarrow$ lectures by dev.